

L13 – Week 7

# Introduction to Multi-armed Bandits (part 2)

CS 295 Optimization for Machine Learning

Ioannis Panageas

# Recap of framework (stochastic)

**Setting.** We are given  $K$  arms and time window  $T$  (known). At each time step  $t = 1 \dots T$ .

- Player chooses arm  $a_t$ .
- Observes reward  $r_t \in [0, 1]$  for the chosen arm.
- The algorithm observes only the reward for the selected action, and nothing else.
- The reward for each action is IID. For each arm  $a \in [K]$ , there is a distribution  $D_a$  over reals, called the reward distribution (**unknown**). Every time this action is chosen, the reward is sampled independently from this distribution.

**Goal:** Minimize the regret

$$R(T) = \mu^* T - \sum_{t=1}^T \mu(a_t) \text{ or } \mathbb{E}[R(T)].$$

# Upper Confidence Bounds

**Definition (Confidence bounds).** We define upper/lower confidence bounds for every arm  $a$  and round  $t$

$$UCB_t(a) = \hat{\mu}_t(a) + r_t(a), \quad LCB_t(a) = \hat{\mu}_t(a) - r_t(a),$$

where  $\hat{\mu}_t(a)$  is the average reward of arm  $a$  so far,  $r_t(a) = \sqrt{\frac{2 \log T}{n_t(a)}}$  (**confidence radius**) and  $n_t(a)$  is the number of samples from arm  $a$  in round  $1, \dots, t$ ,

# Upper Confidence Bounds

**Definition (Confidence bounds).** We define upper/lower confidence bounds for every arm  $a$  and round  $t$

$$UCB_t(a) = \hat{\mu}_t(a) + r_t(a), \quad LCB_t(a) = \hat{\mu}_t(a) - r_t(a),$$

where  $\hat{\mu}_t(a)$  is the average reward of arm  $a$  so far,  $r_t(a) = \sqrt{\frac{2 \log T}{n_t(a)}}$  (**confidence radius**) and  $n_t(a)$  is the number of samples from arm  $a$  in round  $1, \dots, t$ ,

**Definition (UCB).** Consider the following algorithm:

1. **Try** each arm once.
2. In each round  $t$ , **pick**  $\arg \max_a UCB_t(a)$ .

# Analysis of UCB

Remarks:

- An arm  $a$  has the largest  $UCB_t$  for two reasons: The **empirical reward is large** (hence it is likely  $a$  has high reward) or **confidence radius is large**, thus the arm has not been explored much.

**Either reason makes this arm worth choosing!**

# Analysis of UCB

Remarks:

- An arm  $a$  has the largest  $UCB_t$  for two reasons: The **empirical reward is large** (hence it is likely  $a$  has high reward) or **confidence radius is large**, thus the arm has not been explored much.

**Either reason makes this arm worth choosing!**

**Theorem (Regret).** *UCB algorithm achieves regret*

$$\mathbb{E}[R(T)] \text{ to be } O(\sqrt{KT \log T}).$$

**Theorem (Regret v2).** *UCB algorithm achieves regret*

$$\mathbb{E}[R(T)] \leq O(\log T) \left( \sum_{a: \mu(a) < \mu(a^*)} \frac{1}{\mu(a^*) - \mu(a)} \right).$$

# Analysis of UCB

Let us define the “clean” event (we condition on that)

$$\mathcal{E} = \{\forall j, a \mid \hat{\mu}_j(a) - \mu(a) \leq r_j(a)\}.$$

Let  $a^*$  be an optimal arm and assume that we chose arm  $a_t$  at time  $t$  then:

$$\mu(a_t) + 2r_t(a_t) \geq \hat{\mu}_t(a_t) + r_t(a_t) \text{ clean event}$$

# Analysis of UCB

Let us define the “clean” event (we condition on that)

$$\mathcal{E} = \{\forall j, a \mid \hat{\mu}_j(a) - \mu(a) \leq r_j(a)\}.$$

Let  $a^*$  be an optimal arm and assume that we chose arm  $a_t$  at time  $t$  then:

$$\begin{aligned} \mu(a_t) + 2r_t(a_t) &\geq \hat{\mu}_t(a_t) + r_t(a_t) \text{ clean event} \\ &= \text{UCB}_t(a_t) \end{aligned}$$



# Analysis of UCB

Let us define the “clean” event (we condition on that)

$$\mathcal{E} = \{\forall j, a \mid \hat{\mu}_j(a) - \mu(a) \leq r_j(a)\}.$$

Let  $a^*$  be an optimal arm and assume that we chose arm  $a_t$  at time  $t$  then:

$$\begin{aligned} \mu(a_t) + 2r_t(a_t) &\geq \hat{\mu}_t(a_t) + r_t(a_t) \text{ clean event} \\ &= \text{UCB}_t(a_t) \\ &\geq \text{UCB}_t(a^*) \text{ since we chose } a_t \end{aligned}$$

# Analysis of UCB

Let us define the “clean” event (we condition on that)

$$\mathcal{E} = \{\forall j, a \mid \hat{\mu}_j(a) - \mu(a) \leq r_j(a)\}.$$

Let  $a^*$  be an optimal arm and assume that we chose arm  $a_t$  at time  $t$  then:

$$\begin{aligned} \mu(a_t) + 2r_t(a_t) &\geq \hat{\mu}_t(a_t) + r_t(a_t) \text{ clean event} \\ &= \text{UCB}_t(a_t) \\ &\geq \text{UCB}_t(a^*) \text{ since we chose } a_t \\ &= \mu_t(a^*) + r_t(a^*) \geq \mu(a^*) \text{ clean event.} \end{aligned}$$

# Analysis of UCB

Let us define the “clean” event (we condition on that)

$$\mathcal{E} = \{\forall j, a \mid \hat{\mu}_j(a) - \mu(a) \leq r_j(a)\}.$$

Let  $a^*$  be an optimal arm and assume that we chose arm  $a_t$  at time  $t$  then:

$$\begin{aligned} \mu(a_t) + 2r_t(a_t) &\geq \hat{\mu}_t(a_t) + r_t(a_t) \text{ clean event} \\ &= \text{UCB}_t(a_t) \\ &\geq \text{UCB}_t(a^*) \text{ since we chose } a_t \\ &= \mu_t(a^*) + r_t(a^*) \geq \mu(a^*) \text{ clean event.} \end{aligned}$$

Hence it holds

$$2r_t(a_t) \geq \mu(a^*) - \mu(a_t) = \Delta(a_t).$$

# Analysis of UCB

Hence it holds

$$2r_t(a_t) \geq \mu(a^*) - \mu(a_t) = \Delta(a_t).$$

For each arm  $a$  consider the last time  $\tau$  that  $a$  was pulled, then we get  $n_T(a) = n_\tau(a)$  and  $r_T(a) = r_\tau(a)$ . We conclude that

# Analysis of UCB

Hence it holds

$$2r_t(a_t) \geq \mu(a^*) - \mu(a_t) = \Delta(a_t).$$

For each arm  $a$  consider the last time  $\tau$  that  $a$  was pulled, then we get  $n_T(a) = n_\tau(a)$  and  $r_T(a) = r_\tau(a)$ . We conclude that

$$2\sqrt{\frac{2 \log T}{n_T(a)}} = 2r_T(a) \geq \mu(a^*) - \mu(a) = \Delta(a).$$

# Analysis of UCB

Hence it holds

$$2r_t(a_t) \geq \mu(a^*) - \mu(a_t) = \Delta(a_t).$$

For each arm  $a$  consider the last time  $\tau$  that  $a$  was pulled, then we get  $n_T(a) = n_\tau(a)$  and  $r_T(a) = r_\tau(a)$ . We conclude that

$$2\sqrt{\frac{2 \log T}{n_T(a)}} = 2r_T(a) \geq \mu(a^*) - \mu(a) = \Delta(a).$$

The contribution of arm  $a$  to the total regret is

$$\Delta(a) \times n_T(a) \leq 2\sqrt{2n_T(a) \log T}.$$

# Analysis of UCB

Hence it holds

$$2r_t(a_t) \geq \mu(a^*) - \mu(a_t) = \Delta(a_t).$$

For each arm  $a$  consider the last time  $\tau$  that  $a$  was pulled, then we get  $n_T(a) = n_\tau(a)$  and  $r_T(a) = r_\tau(a)$ . We conclude that

$$2\sqrt{\frac{2 \log T}{n_T(a)}} = 2r_T(a) \geq \mu(a^*) - \mu(a) = \Delta(a).$$

The contribution of arm  $a$  to the total regret is

$$\Delta(a) \times n_T(a) \leq 2\sqrt{2n_T(a) \log T}.$$

Hence the regret is bounded by

$$2\sqrt{2 \log T} \sum_a \sqrt{n_T(a)}.$$

# Analysis of UCB

Hence the regret is bounded by

$$2\sqrt{2\log T} \sum_a \sqrt{n_T(a)}.$$

Finally observe that  $\sqrt{x}$  is a concave function hence, by Jensen's inequality we get

$$\frac{1}{K} \sum_a \sqrt{n_T(a)} \leq \sqrt{\frac{1}{K} \sum_a n_T(a)} \leq \sqrt{\frac{T}{K}}.$$



# Analysis of UCB

Hence the regret is bounded by

$$2\sqrt{2\log T} \sum_a \sqrt{n_T(a)}.$$

Finally observe that  $\sqrt{x}$  is a concave function hence, by Jensen's inequality we get

$$\frac{1}{K} \sum_a \sqrt{n_T(a)} \leq \sqrt{\frac{1}{K} \sum_a n_T(a)} \leq \sqrt{\frac{T}{K}}.$$

We conclude that the regret is bounded by

$$O(\sqrt{TK \log T}).$$

# Analysis of UCB

Recall that we showed

$$2\sqrt{\frac{2\log T}{n_T(a)}} = 2r_T(a) \geq \mu(a^*) - \mu(a) = \Delta(a).$$

This implies that

$$n_T(a) \leq \frac{8\log T}{\Delta(a)^2}$$

# Analysis of UCB

Recall that we showed

$$2\sqrt{\frac{2\log T}{n_T(a)}} = 2r_T(a) \geq \mu(a^*) - \mu(a) = \Delta(a).$$

This implies that

$$n_T(a) \leq \frac{8\log T}{\Delta(a)^2}$$

The contribution of arm  $a$  to the total regret is

$$\Delta(a) \times n_T(a) \leq \frac{8\log T}{\Delta(a)}.$$

Hence the regret is bounded by

$$O(\log T) \sum_a \frac{1}{\Delta(a)}.$$

# Framework (adversarial bandits)

**Setting.** We are given  $K$  arms and time window  $T$  (known). At each time step  $t = 1 \dots T$ .

- *Player* chooses arm  $a_t$ .
  - *Adversary* picks cost  $c_t(a)$  for each arm  $a$ .
  - *Player* observes cost  $c_t(a_t) \in [0, 1]$  for the chosen arm.
- 
- The player observes only the cost for the selected action, and nothing else.

**Goal:** Minimize the regret

# Framework (adversarial bandits)

**Setting.** We are given  $K$  arms and time window  $T$  (known). At each time step  $t = 1 \dots T$ .

- *Player* chooses arm  $a_t$ .
  - *Adversary* picks cost  $c_t(a)$  for each arm  $a$ .
  - *Player* observes cost  $c_t(a_t) \in [0, 1]$  for the chosen arm.
- 
- The player observes only the cost for the selected action, and nothing else.

**Goal:** Minimize the regret

$$R(T) = \sum_{t \in [T]} c_t(a_t) - \min_a \sum_{t \in [T]} c_t(a) \text{ or } \mathbb{E}[R(T)].$$

# MWU (recap)

**Algorithm (MWUA).** We define the following algorithm:

1. Initialize  $w_i^0 = 1$  for all  $i \in [n]$ .
2. **For**  $t=1 \dots T$  **do**
3.     **Choose** action  $i$  with probability proportional to  $w_i^{t-1}$ .
4.     **For** each action  $i$  **do**
5.          $w_i^t = (1 - \epsilon)^{c_i^t} w_i^{t-1}$ .
6.     **End For**
7. **End For**

Remarks:

- We choose  $i$  with probability  $p_i^t = \frac{w_i^{t-1}}{\sum_j w_j^{t-1}}$ .
- $c_i^t$  is the cost of action  $i$  at time  $t$  chosen by the adversary.

Can we use this for adversarial bandits? Reduction

# Exp3 Algorithm

**Algorithm (Exp3).** We define the following algorithm:

1. Initialize  $w_i^0 = 1$  for all  $i \in [n]$ .
2. **For**  $t=1 \dots T$  **do**
3.     **Choose** action  $i$  with probability proportional to  $w_i^{t-1}$ .
4.     **Only for** the chosen action (say  $i$ ) **do**
5.          $w_i^t = (1 - \epsilon)^{c_i^t/p_i^t} w_i^{t-1}$ .
6.     **End For**
7. **End For**

Remarks:

- We choose  $i$  with probability  $p_i^t = \frac{w_i^{t-1}}{\sum_j w_j^{t-1}}$ .
- $c_i^t$  is the cost of action  $i$  at time  $t$  chosen by the adversary.
- Essentially, we assume that **all the actions got cost zero except the chosen action that got cost  $\hat{c}_i^t := c_i^t/p_i^t$ .**

# Exp3 Algorithm

**Algorithm (Exp3).** We define the following algorithm:

1. Initialize  $w_i^0 = 1$  for all  $i \in [n]$ .
2. **For**  $t=1 \dots T$  **do**
3.     **Choose** action  $i$  with probability proportional to  $w_i^{t-1}$ .
4.     **Only for** the chosen action (say  $i$ ) **do**
5.          $w_i^t = (1 - \epsilon)^{c_i^t/p_i^t} w_i^{t-1}$ .
6.     **End For**
7. **End For**

Remarks:

- We choose  $i$  with probability  $p_i^t = \frac{w_i^{t-1}}{\sum_j w_j^{t-1}}$ .
- $c_i^t$  is the cost of action  $i$  at time  $t$  chosen by the adversary.
- Essentially, we assume that **all the actions got cost zero except the chosen action that got cost  $\hat{c}_i^t := c_i^t/p_i^t$ .**

**What is the cost of every action? Each a r.v that is an unbiased estimator!**

Formally we ensure that  $\mathbb{E}[\hat{c}_i^t | p^t] = c_i^t$  for all  $i$ .



# Exp3 Algorithm

**Algorithm (Exp3).** We define the following algorithm:

1. Initialize  $w_i^0 = 1$  for all  $i \in [n]$ .
2. **For**  $t=1 \dots T$  **do**
3.     **Choose** action  $i$  with probability proportional to  $w_i^{t-1}$ .
4.     **Only for** the chosen action (say  $i$ ) **do**
5.          $w_i^t = (1 - \epsilon)^{c_i^t/p_i^t} w_i^{t-1}$ .
6.     **End For**
7. **End For**

Remarks:

- We choose  $i$  with probability  $p_i^t = \frac{w_i^{t-1}}{\sum_j w_j^{t-1}}$ .
- $c_i^t$  is the cost of action  $i$  at time  $t$  chosen by the adversary.
- Essentially, we assume that **all the actions got cost zero except the chosen action that got cost  $\hat{c}_i^t := c_i^t/p_i^t$ .**

**What is the cost of every action? Each a r.v that is an unbiased estimator!**

Formally we ensure that  $\mathbb{E}[\hat{c}_i^t | p^t] = c_i^t$  for all  $i$ .

We will choose  $\epsilon = \sqrt{\frac{2 \log K}{TK}}$  and we will get regret  $O(\sqrt{TK \log K})$ .

# Analysis of Exp3

Recall that for the analysis of MWU we defined a potential function  $\Phi_t$  (sum of weights).

$$\text{We set } \Phi_t = -\frac{1}{\epsilon} \log \sum_i e^{-\epsilon \sum_{\tau=1}^{t-1} \hat{c}_i^\tau}.$$

Set  $L_i^t = \sum_{\tau=1}^t \hat{c}_i^\tau$ . Observe that

# Analysis of Exp3

Recall that for the analysis of MWU we defined a potential function  $\Phi_t$  (sum of weights).

$$\text{We set } \Phi_t = -\frac{1}{\epsilon} \log \sum_i e^{-\epsilon \sum_{\tau=1}^{t-1} \hat{c}_i^\tau}.$$

Set  $L_i^t = \sum_{\tau=1}^t \hat{c}_i^\tau$ . Observe that

$$\Phi_{t+1} - \Phi_t = -\frac{1}{\epsilon} \log \frac{\sum_i e^{-\epsilon L_i^t}}{\sum_i e^{-\epsilon L_i^{t-1}}}$$

# Analysis of Exp3

Recall that for the analysis of MWU we defined a potential function  $\Phi_t$  (sum of weights).

$$\text{We set } \Phi_t = -\frac{1}{\epsilon} \log \sum_i e^{-\epsilon \sum_{\tau=1}^{t-1} \hat{c}_i^\tau}.$$

Set  $L_i^t = \sum_{\tau=1}^t \hat{c}_i^\tau$ . Observe that

$$\begin{aligned} \Phi_{t+1} - \Phi_t &= -\frac{1}{\epsilon} \log \frac{\sum_i e^{-\epsilon L_i^t}}{\sum_i e^{-\epsilon L_i^{t-1}}} \\ &= -\frac{1}{\epsilon} \log \frac{\sum_i e^{-\epsilon L_i^{t-1}} e^{-\epsilon \hat{c}_i^t}}{\sum_i e^{-\epsilon L_i^{t-1}}} \end{aligned}$$

# Analysis of Exp3

Recall that for the analysis of MWU we defined a potential function  $\Phi_t$  (sum of weights).

$$\text{We set } \Phi_t = -\frac{1}{\epsilon} \log \sum_i e^{-\epsilon \sum_{\tau=1}^{t-1} \hat{c}_i^\tau}.$$

Set  $L_i^t = \sum_{\tau=1}^t \hat{c}_i^\tau$ . Observe that

$$\begin{aligned} \Phi_{t+1} - \Phi_t &= -\frac{1}{\epsilon} \log \frac{\sum_i e^{-\epsilon L_i^t}}{\sum_i e^{-\epsilon L_i^{t-1}}} \\ &= -\frac{1}{\epsilon} \log \frac{\sum_i e^{-\epsilon L_i^{t-1}} e^{-\epsilon \hat{c}_i^t}}{\sum_i e^{-\epsilon L_i^{t-1}}} \\ &= -\frac{1}{\epsilon} \log \mathbb{E}_{i \sim p^t} [e^{-\epsilon \hat{c}_i^t}] \end{aligned}$$

# Analysis of Exp3

Recall that for the analysis of MWU we defined a potential function  $\Phi_t$  (sum of weights).

$$\text{We set } \Phi_t = -\frac{1}{\epsilon} \log \sum_i e^{-\epsilon \sum_{\tau=1}^{t-1} \hat{c}_i^\tau}.$$

Set  $L_i^t = \sum_{\tau=1}^t \hat{c}_i^\tau$ . Observe that

$$\begin{aligned} \Phi_{t+1} - \Phi_t &= -\frac{1}{\epsilon} \log \frac{\sum_i e^{-\epsilon L_i^t}}{\sum_i e^{-\epsilon L_i^{t-1}}} \\ &= -\frac{1}{\epsilon} \log \frac{\sum_i e^{-\epsilon L_i^{t-1}} e^{-\epsilon \hat{c}_i^t}}{\sum_i e^{-\epsilon L_i^{t-1}}} \\ &= -\frac{1}{\epsilon} \log \mathbb{E}_{i \sim p^t} [e^{-\epsilon \hat{c}_i^t}] \\ &\geq -\frac{1}{\epsilon} \log \mathbb{E}_{i \sim p^t} [1 - \epsilon \hat{c}_i^t + \frac{1}{2} \epsilon^2 \hat{c}_i^t{}^2] \end{aligned}$$

Since  $e^{-x} \leq 1 - x + \frac{1}{2}x^2$ .

# Analysis of Exp3

$$\begin{aligned}\Phi_{t+1} - \Phi_t &\geq -\frac{1}{\epsilon} \log \mathbb{E}_{i \sim p^t} \left[ 1 - \epsilon \hat{c}_i^t + \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \\ &= -\frac{1}{\epsilon} \log \left( 1 - \mathbb{E}_{i \sim p^t} \left[ \epsilon \hat{c}_i^t - \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \right)\end{aligned}$$

# Analysis of Exp3

$$\begin{aligned}\Phi_{t+1} - \Phi_t &\geq -\frac{1}{\epsilon} \log \mathbb{E}_{i \sim p^t} \left[ 1 - \epsilon \hat{c}_i^t + \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \\ &= -\frac{1}{\epsilon} \log \left( 1 - \mathbb{E}_{i \sim p^t} \left[ \epsilon \hat{c}_i^t - \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \right) \\ &\geq \frac{1}{\epsilon} \mathbb{E}_{i \sim p^t} \left[ \epsilon \hat{c}_i^t - \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right]\end{aligned}$$



# Analysis of Exp3

$$\begin{aligned}\Phi_{t+1} - \Phi_t &\geq -\frac{1}{\epsilon} \log \mathbb{E}_{i \sim p^t} \left[ 1 - \epsilon \hat{c}_i^t + \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \\ &= -\frac{1}{\epsilon} \log \left( 1 - \mathbb{E}_{i \sim p^t} \left[ \epsilon \hat{c}_i^t - \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \right) \\ &\geq \frac{1}{\epsilon} \mathbb{E}_{i \sim p^t} \left[ \epsilon \hat{c}_i^t - \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \\ &= \sum_i p_i^t \hat{c}_i^t - \frac{1}{2} \epsilon \sum_i p_i^t \hat{c}_i^{t^2}\end{aligned}$$

# Analysis of Exp3

$$\begin{aligned}\Phi_{t+1} - \Phi_t &\geq -\frac{1}{\epsilon} \log \mathbb{E}_{i \sim p^t} \left[ 1 - \epsilon \hat{c}_i^t + \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \\ &= -\frac{1}{\epsilon} \log \left( 1 - \mathbb{E}_{i \sim p^t} \left[ \epsilon \hat{c}_i^t - \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \right) \\ &\geq \frac{1}{\epsilon} \mathbb{E}_{i \sim p^t} \left[ \epsilon \hat{c}_i^t - \frac{1}{2} \epsilon^2 \hat{c}_i^{t^2} \right] \\ &= \sum_i p_i^t \hat{c}_i^t - \frac{1}{2} \epsilon \sum_i p_i^t \hat{c}_i^{t^2}\end{aligned}$$

By taking expectation we get

$$\mathbb{E}[\Phi_{t+1} - \Phi_t] \geq \sum_i p_i^t c_i^t - \frac{1}{2} \epsilon \sum_i c_i^{t^2} \geq \sum_i p_i^t c_i^t - \frac{K\epsilon}{2}$$

# Analysis of Exp3

We conclude that (telescopic sum)

$$\mathbb{E}[\Phi_T - \Phi_1] \geq \sum_{t=1}^T \sum_i p_i^t c_i^t - \frac{KT\epsilon}{2}$$

# Analysis of Exp3

We conclude that (telescopic sum)

$$\mathbb{E}[\Phi_T - \Phi_1] \geq \sum_{t=1}^T \sum_i p_i^t c_i^t - \frac{KT\epsilon}{2}$$

Finally

$$\mathbb{E}[\Phi_T - \Phi_1] \leq \mathbb{E}[L_{i^*}^T - (-\frac{1}{\epsilon} \log K)] = \sum_t c_{i^*}^t + \frac{1}{\epsilon} \log K.$$

Hence

$$\mathbb{E}[R(T)] = \sum_t \sum_i p_i^t c_i^t - \sum_t c_{i^*}^t \leq \frac{KT\epsilon}{2} + \frac{1}{\epsilon} \log K$$

# Analysis of Exp3

We conclude that (telescopic sum)

$$\mathbb{E}[\Phi_T - \Phi_1] \geq \sum_{t=1}^T \sum_i p_i^t c_i^t - \frac{KT\epsilon}{2}$$

Finally

$$\mathbb{E}[\Phi_T - \Phi_1] \leq \mathbb{E}[L_{i^*}^T - (-\frac{1}{\epsilon} \log K)] = \sum_t c_{i^*}^t + \frac{1}{\epsilon} \log K.$$

Hence

$$\mathbb{E}[R(T)] = \sum_t \sum_i p_i^t c_i^t - \sum_t c_{i^*}^t \leq \frac{KT\epsilon}{2} + \frac{1}{\epsilon} \log K$$

We choose  $\epsilon = \sqrt{\frac{2 \log K}{TK}}$  and it follows that  $\mathbb{E}[R(T)]$  is  $O(\sqrt{TK \log K})$ .

# Conclusion

- Introduction to Multi-armed bandits.
  - UCB.
  - Exp3
- Next lecture we will talk about basics in **Markov Decision Processes and Stochastic Games.**